

Bakery Algorithm

This algorithm solves the critical section problem for n processes in software. The basic idea is that of a bakery; customers take numbers, and whoever has the lowest number gets service next. Here, of course, “service” means entry to the critical section.

```
1 var choosing: shared array[0..n-1] of boolean;
2   number: shared array[0..n-1] of integer;
3   ...
4 repeat
5   choosing[i] := true;
6   number[i] := max(number[0], number[1], ..., number[n-1]) + 1;
7   choosing[i] := false;
8   for j := 0 to n-1 do begin
9     while choosing[j] do (* nothing *);
10    while number[j] <> 0 and number[j], j) < (number[i], i) do
11      (* nothing *);
12    end;
13    (* critical section *)
14    number[i] := 0;
15    (* remainder section *)
16 until false;
```

lines 1–2: Here, `choosing[i]` is true if process i is choosing a number. The number that process i will use to enter the critical section is in `number[i]`; it is 0 if process i is not trying to enter its critical section.

lines 4–6: These three lines first indicate that the process is choosing a number (line 4), then try to assign a unique number to the process process i (line 5); however, that does not always happen. Afterwards, process i indicates it is done (line 6).

lines 8–11: Now we select which process goes into the critical section. Process i waits until it has the lowest number of all the processes waiting to enter the critical section. If two processes have the same number, the one with the smaller name — the value of the index — goes in; the notation “ $(a,b) < (c,d)$ ”; means true if $a < c$ or if both $a = c$ and $b < d$ (lines 9–10). Note that if a process is not trying to enter the critical section, its number is 0. Also, if a process is choosing a number when process i tries to look at it, process i waits until it has done so before looking (line 8).

line 14: Now process i is no longer interested in entering its critical section, so it sets `number[i]` to 0.