

## Outline for April 21, 2005

1. Policy
  - a. Policy languages: high level, low level
2. Bell-LaPadula Model (security classifications only)
  - a. Go through security clearance, classification
  - b. Describe simple security condition (no reads up), \*-property (no writes down), discretionary security property
  - c. State Basic Security Theorem: if it's secure and transformations follow these rules, it's still secure
3. Bell-LaPadula Model (security levels)
  - a. Go through security clearance, categories, levels
4. Lattice models
  - a. Poset,  $\leq$  the relation
  - b. Reflexive, antisymmetric, transitive
  - c. Greatest lower bound, least upper bound
  - d. Example with complex numbers
5. Bell-LaPadula Model
  - a. Apply lattice work
    - i. Set of classes SC is a partially ordered set under relation  $\leq$  with GLB (greatest lower bound), LUB (least upper bound) operators
    - ii. Note: is reflexive, transitive, antisymmetric
    - iii. Examples:  $(A, C) \leq (A', C')$  iff  $A \leq A'$  and  $C \subseteq C'$ ;  
 $LUB((A, C), (A', C')) = (\max(A, A'), C \cup C')$ ,  $GLB((A, C), (A', C')) = (\min(A, A'), C \cap C')$
  - b. Describe simple security condition (no reads up), \*-property (no writes down), discretionary security property
  - c. State Basic Security Theorem: if it's secure and transformations follow these rules, it's still secure
  - d. Maximum, current security level
6. Example: DG/UX UNIX
  - a. Labels and regions
  - b. Multilevel directories
  - c. File object labels
  - d. MAC tuples
7. BLP: formally
  - a. Elements of system:  $s_i$  subjects,  $o_i$  objects
  - b. State space  $V = B \times M \times F \times H$  where:
    - $B$  set of current accesses (i.e., access modes each subject has currently to each object);
    - $M$  access permission matrix;
    - $F$  consists of 3 functions:  $f_s$  is security level associated with each subject,  $f_o$  security level associated with each object, and  $f_c$  current security level for each subject
    - $H$  hierarchy of system objects, functions  $h: O \rightarrow P(O)$  with two properties:  
 If  $o_i \neq o_j$ , then  $h(o_i) \cap h(o_j) = \emptyset$   
 There is no set  $\{ o_1, \dots, o_k \} \subseteq O$  such that for each  $i$ ,  $o_{i+1} \in h(o_i)$  and  $o_{k+1} = o_1$ .
  - c. Set of requests is  $R$
  - d. Set of decisions is  $D$
  - e.  $W \subseteq R \times D \times V \times V$  is motion from one state to another.
  - f. System  $\Sigma(R, D, W, z_0) \subseteq X \times Y \times Z$  such that  $(x, y, z) \in \Sigma(R, D, W, z_0)$  iff  $(x_t, y_t, z_t, z_{t-1}) \in W$  for each  $i \in T$ ; latter is an action of system
  - g. Theorem:  $\Sigma(R, D, W, z_0)$  satisfies the simple security property for any initial state  $z_0$  that satisfies the simple security property iff  $W$  satisfies the following conditions for each action  $(r_i, d_i, (b', m', f', h'), (b, m, f, h))$ :

- i. each  $(s, o, x) \in b' - b$  satisfies the simple security condition relative to  $f'$  (i.e.,  $x$  is not read, or  $x$  is read and  $f'_s(s)$  dominates  $f_o(o)$ )
  - ii. if  $(s, o, x) \in b$  does not satisfy the simple security condition relative to  $f'$ , then  $(s, o, x) \notin b'$
  - h. Theorem:  $\Sigma(R, D, W, z_0)$  satisfies the \*-property relative to  $S' \subseteq S$ , for any initial state  $z_0$  that satisfies the \*-property relative to  $S'$  iff  $W$  satisfies the following conditions for each  $(r_i, d_i, (b', m', f', h'), (b, m, f, h))$ :
    - i. for each  $s \in S'$ , any  $(s, o, x) \in b' - b$  satisfies the \*-property with respect to  $f'$
    - ii. for each  $s \in S'$ , if  $(s, o, x) \in b$  does not satisfy the \*-property with respect to  $f'$ , then  $(s, o, x) \notin b'$
  - i. Theorem:  $\Sigma(R, D, W, z_0)$  satisfies the ds-property iff the initial state  $z_0$  satisfies the ds-property and  $W$  satisfies the following conditions for each action  $(r_i, d_i, (b', m', f', h'), (b, m, f, h))$ :
    - i. if  $(s, o, x) \in b' - b$ , then  $x \in m'[s, o]$ ;
    - ii. if  $(s, o, x) \in b$  and  $x \in m'[s, o]$  then  $(s, o, x) \notin b'$
  - j. Basic Security Theorem: A system  $\Sigma(R, D, W, z_0)$  is secure iff  $z_0$  is a secure state and  $W$  satisfies the conditions of the above three theorems for each action.
8. BLP: formally
- a. Define ssc-preserving, \*-property-preserving, ds-property-preserving
  - b. Define relation  $W(\omega)$
  - c. Show conditions under which rules are ssc-preserving, \*-property-preserving, ds-property-preserving
  - d. Show when adding a state preserves those properties
  - e. Example instantiation: *get-read* for Multics
9. Tranquility
- a. Strong tranquility
  - b. Weak tranquility
10. System Z and the controversy