

ECS 235B Module 4

Access Control Matrix

Description

objects (entities)

| | O_1 | ... | O_m | S_1 | ... | S_n |
|-------|-------|-----|-------|-------|-----|-------|
| s_1 | | | | | | |
| s_2 | | | | | | |
| ... | | | | | | |
| s_n | | | | | | |

subjects

- Subjects $S = \{ s_1, \dots, s_n \}$
- Objects $O = \{ o_1, \dots, o_m \}$
- Rights $R = \{ r_1, \dots, r_k \}$
- Entries $A[s_i, o_j] \subseteq R$
- $A[s_i, o_j] = \{ r_x, \dots, r_y \}$ means subject s_i has rights r_x, \dots, r_y over object o_j

Example 1

- Processes p, q
- Files f, g
- Rights r, w, x, a, o

| | f | g | p | q |
|-----|-------|------|--------|--------|
| p | rwo | r | $rwxo$ | w |
| q | a | ro | r | $rwxo$ |

Example 2

- Host names *telegraph*, *nob*, *toadflax*
- Rights *own*, *ftp*, *nfs*, *mail*

| | <i>telegraph</i> | <i>nob</i> | <i>toadflax</i> |
|------------------|------------------|----------------------------|----------------------------|
| <i>telegraph</i> | <i>own</i> | <i>ftp</i> | <i>ftp</i> |
| <i>nob</i> | | <i>ftp, mail, nfs, own</i> | <i>ftp, nfs, mail</i> |
| <i>toadflax</i> | | <i>ftp, mail</i> | <i>ftp, mail, nfs, own</i> |

Example 3

- Procedures *inc_ctr*, *dec_ctr*, *manage*
- Variable *counter*
- Rights *+*, *-*, *call*

| | <i>counter</i> | <i>inc_ctr</i> | <i>dec_ctr</i> | <i>manage</i> |
|----------------|----------------|----------------|----------------|---------------|
| <i>inc_ctr</i> | <i>+</i> | | | |
| <i>dec_ctr</i> | <i>-</i> | | | |
| <i>manager</i> | | <i>call</i> | <i>call</i> | <i>call</i> |

UNIX/Linux Access Controls

- Files

- A is `~/bishop/a.out` (0755, or `rwxr-xr-x`)
- B is `/etc/passwd` (0644, or `rw-r--r--`)
- H is `/home/bishop` (0711, or `rwX--X--X`)
- S is `/bin/su` (4711, or `s--rwx--X--X`)

| | <i>A</i> | <i>B</i> | <i>S</i> | <i>H</i> |
|---------------|-------------|------------|-------------|-------------|
| <i>bishop</i> | <i>rwXO</i> | <i>r</i> | <i>X</i> | <i>rwXO</i> |
| <i>zheng</i> | <i>rx</i> | <i>r</i> | <i>X</i> | <i>X</i> |
| <i>root</i> | <i>rwX</i> | <i>rWO</i> | <i>rwXO</i> | <i>rwX</i> |

UNIX/Linux Access Controls

- Access control matrices are dynamic:
- After bishop executes `chmod 700 /home/bishop`:
 - Same as `chmod u=rwx,g-rwx,o-rwx /home/bishop`

| | <i>A</i> | <i>B</i> | <i>S</i> | <i>H</i> |
|---------------|-------------|------------|-------------|-------------|
| <i>bishop</i> | <i>rwxo</i> | <i>r</i> | <i>x</i> | <i>rwxo</i> |
| <i>muwei</i> | | <i>rx</i> | | |
| <i>root</i> | <i>rwX</i> | <i>rwo</i> | <i>rwXO</i> | <i>rwX</i> |

Boolean Expression Evaluation

- ACM controls access to database fields
 - Subjects have attributes
 - Verbs define type of access
 - Rules associated with objects, verb pair
- Subject attempts to access object
 - Rule for object, verb evaluated, grants or denies access

Example

- Subject annie
 - Attributes *role* (artist), *group* (creative)
- Verb paint
 - Default 0 (deny unless explicitly granted)
- Object picture
 - Rule:
paint: 'artist' in subject.role and
'creative' in subject.groups and
time.hour ≥ 0 and time.hour ≤ 4

ACM at 3AM and 10AM

At 3AM, time condition met
ACM is:

... picture ...

| | | | |
|---------------|-------|--|--|
| ... annie ... | | | |
| | paint | | |
| | | | |

At 10AM, time condition not met
ACM is:

... picture ...

| | | | |
|---------------|--|--|--|
| ... annie ... | | | |
| | | | |
| | | | |

History

- Problem: what a process has accessed may affect what it can access now
- Example: procedure in a web applet can access other procedures depending on what procedures it has already accessed
 - S set of *static rights* associated with procedure
 - C set of *current rights* associated with each executing process
 - When process calls procedure, rights are $S \cap C$

Example Program

```
// This routine has no filesystem access rights  
// beyond those in a limited, temporary area
```

```
procedure helper_proc()  
    return sys_kernel_file
```

```
// But this has the right to delete files
```

```
program main()  
    sys_load_file(helper_proc)  
    tmp_file = helper_proc()  
    sys_delete_file(tmp_file)
```

- *sys_kernel_file* contains system kernel
- *tmp_file* is in limited area that *helper_proc()* can access

Before *helper_proc* Called

- Static rights of program

| | <i>sys_kernel_file</i> | <i>tmp_file</i> |
|--------------------|------------------------|-----------------|
| <i>main</i> | delete | delete |
| <i>helper_proc</i> | | delete |

- When program starts, current rights:

| | <i>sys_kernel_file</i> | <i>tmp_file</i> |
|--------------------|------------------------|-----------------|
| <i>main</i> | delete | delete |
| <i>helper_proc</i> | | delete |
| <i>process</i> | delete | delete |

After *helper_proc* Called

- Process rights are intersection of static, previous “current” rights:

| | <i>sys_kernel_file</i> | <i>tmp_file</i> |
|--------------------|------------------------|-----------------|
| <i>main</i> | delete | delete |
| <i>helper_proc</i> | | delete |
| <i>process</i> | | delete |

State Transitions

- Change the protection state of system
- \vdash represents transition
 - $X_i \vdash_{\tau} X_{i+1}$: command τ moves system from state X_i to X_{i+1}
 - $X_i \vdash^* Y$: a sequence of commands moves system from state X_i to Y
- Commands often called *transformation procedures*

Primitive Operations

- **create subject s ; create object o**
 - Creates new row, column in ACM; creates new column in ACM
- **destroy subject s ; destroy object o**
 - Deletes row, column from ACM; deletes column from ACM
- **enter r into $A[s, o]$**
 - Adds r rights for subject s over object o
- **delete r from $A[s, o]$**
 - Removes r rights from subject s over object o

Create Subject

- Precondition: $s \notin S$
- Primitive command: **create subject s**
- Postconditions:
 - $S' = S \cup \{s\}, O' = O \cup \{s\}$
 - $(\forall y \in O') [A'[s, y] = \emptyset], (\forall x \in S') [A'[x, s] = \emptyset]$
 - $(\forall x \in S)(\forall y \in O) [A'[x, y] = A[x, y]]$

Create Object

- Precondition: $o \notin O$
- Primitive command: **create object** o
- Postconditions:
 - $S' = S, O' = O \cup \{o\}$
 - $(\forall x \in S') [A'[x, o] = \emptyset]$
 - $(\forall x \in S)(\forall y \in O) [A'[x, y] = A[x, y]]$

Add Right

- Precondition: $s \in S, o \in O$
- Primitive command: **enter r into $A[s, o]$**
- Postconditions:
 - $S' = S, O' = O$
 - $A'[s, o] = A[s, o] \cup \{r\}$
 - $(\forall x \in S')(\forall y \in O' - \{o\}) [A'[x, y] = A[x, y]]$
 - $(\forall x \in S' - \{s\})(\forall y \in O') [A'[x, y] = A[x, y]]$

Delete Right

- Precondition: $s \in S, o \in O$
- Primitive command: **delete r from $A[s, o]$**
- Postconditions:
 - $S' = S, O' = O$
 - $A'[s, o] = A[s, o] - \{r\}$
 - $(\forall x \in S')(\forall y \in O' - \{o\}) [A'[x, y] = A[x, y]]$
 - $(\forall x \in S' - \{s\})(\forall y \in O') [A'[x, y] = A[x, y]]$

Destroy Subject

- Precondition: $s \in S$
- Primitive command: **destroy subject s**
- Postconditions:
 - $S' = S - \{s\}, O' = O - \{s\}$
 - $(\forall y \in O') [A'[s, y] = \emptyset], (\forall x \in S') [A'[x, s] = \emptyset]$
 - $(\forall x \in S')(\forall y \in O') [A'[x, y] = A[x, y]]$

Destroy Object

- Precondition: $o \in O$
- Primitive command: **destroy object o**
- Postconditions:
 - $S' = S, O' = O - \{ o \}$
 - $(\forall x \in S') [A'[x, o] = \emptyset]$
 - $(\forall x \in S')(\forall y \in O') [A'[x, y] = A[x, y]]$

Creating File

- Process p creates file f with r and w permission

```
command create•file( $p$ ,  $f$ )  
    create object  $f$ ;  
    enter own into  $A[p, f]$ ;  
    enter  $r$  into  $A[p, f]$ ;  
    enter  $w$  into  $A[p, f]$ ;  
end
```

Mono-Operational Commands

- Make process p the owner of file g

```
command make-owner( $p, g$ )  
    enter own into  $A[p, g]$ ;  
end
```

- Mono-operational command
 - Single primitive operation in this command

Conditional Commands

- Let p give q r rights over f , if p owns f
command $grant \cdot read \cdot file \cdot 1(p, f, q)$
 if own **in** $A[p, f]$
 then
 enter r **into** $A[q, f];$
 end
- Mono-conditional command
 - Single condition in this command

Biconditional Commands (and)

- Let p give q r and w rights over f , if p owns f and p has c rights over q

```
command grant•read•file•2( $p, f, q$ )  
  if own in  $A[p, f]$  and  $c$  in  $A[p, q]$   
  then  
    enter  $r$  into  $A[q, f]$  ;  
    enter  $w$  into  $A[q, f]$  ;  
end
```

There Is No “or”

- Let p give q r and w rights over f , if p owns f or p has c rights over q

```
command grant•read•file•3( $p, f, q$ )
  if own in  $A[p, f]$ 
  then
    enter  $r$  into  $A[q, f]$ ;
    enter  $w$  into  $A[q, f]$ ;
  end
command grant•read•file•4( $p, f, q$ )
  if  $c$  in  $A[p, q]$ 
  then
    enter  $r$  into  $A[q, f]$ ;
    enter  $w$  into  $A[q, f]$ ;
  end
grant•read•file•3( $p, f, q$ );
grant•read•file•4( $p, f, q$ )
```

General Form

```
command name of command(parameters)  
    if conditions (if many, separate with and)  
    then  
        list of commands or primitive  
        operations to be executed;  
end
```

- Only one **if**, and it must come *before* any primitive operations or subcommands
- When there is an **if**, no commands may follow it (but there can be commands in the body of the **if**)
- There is no **else**

Copy Flag and Right

- Allows possessor to give rights to another
- Often attached to a right (called a *flag*), so only applies to that right
 - r is read right that cannot be copied
 - rc is read right that can be copied
- Is copy flag copied when giving r rights?
 - Depends on model, instantiation of model

Own Right

- Usually allows possessor to change entries in ACM column
 - So owner of object can add, delete rights for others
 - May depend on what system allows
 - Can't give rights to specific (set of) users
 - Can't pass copy flag to specific (set of) users

Attenuation of Privilege

- Principle says you can't increase your rights, or give rights you do not possess
 - Restricts addition of rights within a system
 - Usually *ignored* for owner
 - Why? Owner gives herself rights, gives them to others, deletes her rights.