

ECS 235B Module 21

The Controversy and System Z

Controversy

- McLean:
 - “value of the BST is much overrated since there is a great deal more to security than it captures. Further, what is captured by the BST is so trivial that it is hard to imagine a realistic security model for which it does not hold.”
 - Basis: given assumptions known to be non-secure, BST can prove a non-secure system to be secure

†-Property

- State (b, m, f, h) satisfies the †-property iff for each $s \in S$ the following hold:
 1. $b(s: \underline{a}) \neq \emptyset \Rightarrow [\forall o \in b(s: \underline{a}) [f_c(s) \text{ dom } f_o(o)]]$
 2. $b(s: \underline{w}) \neq \emptyset \Rightarrow [\forall o \in b(s: \underline{w}) [f_o(o) = f_c(s)]]$
 3. $b(s: \underline{r}) \neq \emptyset \Rightarrow [\forall o \in b(s: \underline{r}) [f_c(s) \text{ dom } f_o(o)]]$
- Idea: for writing, subject dominates object; for reading, subject also dominates object
- Differs from *-property in that the mandatory condition for writing is reversed
 - For *-property, it's object dominates subject

Analogues

The following two theorems can be proved

- $\Sigma(R, D, W, z_0)$ satisfies the \dagger -property relative to $S' \subseteq S$ for any secure state z_0 iff for every action $(r, d, (b, m, f, h), (b', m', f', h'))$, W satisfies the following for every $s \in S'$
 - Every $(s, o, p) \in b - b'$ satisfies the \dagger -property relative to S'
 - Every $(s, o, p) \in b'$ that does not satisfy the \dagger -property relative to S' is not in b
- $\Sigma(R, D, W, z_0)$ is a secure system if z_0 is a secure state and W satisfies the conditions for the simple security condition, the \dagger -property, and the ds-property.

Problem

- This system is *clearly* non-secure!
 - Information flows from higher to lower because of the \dagger -property

Discussion

- Role of Basic Security Theorem is to demonstrate that rules preserve security
- Key question: what is security?
 - Bell-LaPadula defines it in terms of 3 properties (simple security condition, *-property, discretionary security property)
 - Theorems are assertions about these properties
 - Rules describe changes to a *particular* system instantiating the model
 - Showing system is secure requires proving rules preserve these 3 properties

Rules and Model

- Nature of rules is irrelevant to model
- Model treats “security” as axiomatic
- Policy defines “security”
 - This instantiates the model
 - Policy reflects the requirements of the systems
- McLean’s definition differs from Bell-LaPadula
 - ... and is not suitable for a confidentiality policy
- Analysts cannot prove “security” definition is appropriate through the model

System Z

- System supporting weak tranquility
- On *any* request, system downgrades *all* subjects and objects to lowest level and adds the requested access permission
 - Let initial state satisfy all 3 properties
 - Successive states also satisfy all 3 properties
- Clearly not secure
 - On first request, everyone can read everything

Reformulation of Secure Action

- Given state that satisfies the 3 properties, the action transforms the system into a state that satisfies these properties and eliminates any accesses present in the transformed state that would violate the property in the initial state, then the action is secure
- BST holds with these modified versions of the 3 properties

Reconsider System Z

- Initial state:
 - subject s , object o
 - $C = \{\text{High}, \text{Low}\}$, $K = \{\text{All}\}$
- Take:
 - $f_c(s) = (\text{Low}, \{\text{All}\})$, $f_o(o) = (\text{High}, \{\text{All}\})$
 - $m[s, o] = \{ \underline{w} \}$, and $b = \{ (s, o, \underline{w}) \}$.
- s requests \underline{r} access to o
- Now:
 - $f'_o(o) = (\text{Low}, \{\text{All}\})$
 - $(s, o, \underline{r}) \in b'$, $m'[s, o] = \{ \underline{r}, \underline{w} \}$

Non-Secure System Z

- As $(s, o, \underline{r}) \in b' - b$ and $f_o(o) \text{ dom } f_c(s)$, access added that was illegal in previous state
 - Under the new version of the Basic Security Theorem, System Z is not secure
 - Under the old version of the Basic Security Theorem, as $f'_c(s) = f'_o(o)$, System Z is secure

Response: What Is Modeling?

- Two types of models
 1. Abstract physical phenomenon to fundamental properties
 2. Begin with axioms and construct a structure to examine the effects of those axioms
- Bell-LaPadula Model developed as a model in the first sense
 - McLean assumes it was developed as a model in the second sense

Reconciling System Z

- Different definitions of security create different results
 - Under one (original definition in Bell-LaPadula Model), System Z is secure
 - Under other (McLean's definition), System Z is not secure