Outline for April 25, 2000

1. What is a cryptosystem?

- a. (M, C, K, D, E)
- b. attacks: known ciphertext, known plaintext, chosen plaintext
- 2. Transposition ciphers
 - a. Show rail-fence cipher as example
 - b. Show anagramming
- 3. Simple substitution ciphers
 - a. Do Cæsar cipher
 - b. Present Vigenère tableau
 - c. Discuss breaking it (Kasiski method).
 - d. Go through one-time pads
- 4. DES
 - a. Product cipher with 64 bits in, 64 bits out, and 16 48-bit round keys generated from 56 bit key
 - b. Note S-boxes are real heart of algorithm
 - c. Complementation property: $DES_k(m) = (DES_{k'}(m'))'$ where x' is the bitwise complement of x;
 - d. Differential cryptanalysis: first version unusable as at 16 rounds, more plaintext/ciphertext pairs needed than exhaustive key trial; but for 15 rounds, cuts this time. Later versions cut it to 2⁴⁷ tries. Works by comparing xors of results with xors of corresponding plaintext.. Designers of DES knew about this one, hence the design of the S-boxes
 - e. Linear cryptanalysis drops required chosen plaintext/ciphertext pairs to 2⁴²; not known to designers of DES.
 - f. Triple DES and EDE mode
- 5. Public Key
 - a. based on NP-hard problems (knapsack)
 - b. based on hard mathematical problems (like factoring)
- 6. Do RSA
 - a. Exponentiation cipher: $C = m^e \mod n$, $M = C^d \mod n$; *d* is private key, (e, n) is public key; must choose *d* first, then *e* so that *ed* mod $\phi(n) = 1$.
 - b. Why? as $ed \mod \phi(n) = 1$, $ed = t\phi(n) + 1$ for some integer *t*. Then
 - $C^d \mod n = (M^e \mod n)^d \mod n$
 - $= M^{ed} \mod n$
 - $= M^{t\phi(n) + 1} \mod n$
 - $= M(M^{t\phi(n)} \mod n) \mod n$
 - $= \mathbf{M}(M^{\phi(n)} \mod n)^t \mod n$
 - $= M(1)^t \mod n$
 - $= M \mod n$

by Fermat's Little Theorem

c. Example: p = 5, q = 7, n = 35, $\phi(n) = 24$; choose e = 11, then d = 11. HELLO WORLD is 07 04 11 11 14 22 14 17 11 03; enciphering is $C = 07^{11} \mod 35 = 28$, *etc.* so encipherment is 28 09 16 16 14 08 14 33 16 12.